

# Brute Force Algorithm

## Algorithm

2014 Fall Semester



# Brute Force

- A straightforward approach, usually based directly on the problem's statement and definitions of the concepts involved
- Examples:
  - Computing  $a^n$  ( $a > 0$ ,  $n$  a nonnegative integer)
  - Computing  $n!$
  - Multiplying two matrices
  - Searching for a key of a given value in a list



# Concepts of Brute Force Algorithm

- A brute-force algorithm solves a problem in the most **simple**, direct or obvious way. As a result, such an algorithm can end up doing far **more work** to solve a given problem than a more clever or sophisticated algorithm might do. On the other hand, a brute-force algorithm is often **easier to implement** than a more sophisticated one and, because of this simplicity, sometimes it can be more efficient.
- Typically, a brute-force algorithm solves such a problem by **exhaustively enumerating all the possibilities**. I.e., for every decision we consider each possible outcome.



# The Simplest Approach

- Brute Force - the simplest of the design strategies
  - Is a straightforward approach to solving a problem, usually directly based on the problem's statement and definitions of the concepts involved.
  - Just do it - the brute-force strategy is easiest to apply.
  - Results in an algorithm that can be improved with a modest amount of time.
- Brute force is important due to its wide applicability and simplicity.

# Examples of Brute Force Algorithm

- Selection Sort
- Bubble Sort
- String Matching
- Sequential Search
- Traveling Salesman Problem (Hamilton Circuits)
- Knapsack Problem
- Job Assignment Problem

# Selection Sort

- Find its smallest element by scanning the list
- Exchange it with the first element making the smallest element in its final position in the sorted list.
- Scan the list, starting with the second element to find the smallest among the last  $n-1$ .
- The second element will be put in its final position.

# Selection Sort – how to work

| 89 , 45 , 68 , 90 , 29 , 34 , **17**

17 | 45 , 68 , 90 , **29** , 34 , 89

17 , 29 | 68 , 90 , 45 , **34** , 89

17 , 29 , 34 | 90 , **45** , 68 , 89

17 , 29 , 34 , 45 | 90 , **68** , 89

17 , 29 , 34 , 45 , 68 | 90 , **89**

17 , 29 , 34 , 45 , 68 , 89 | 90



# Bubble Sort

- Compare adjacent elements of the list.
- Exchange them if they are out of order.
- By doing it repeatedly, we end up “Bubbling up” the largest element to the last position on the list
- The next pass bubbles up the second largest element, and so on until, after  $n-1$  passes, the list is sorted.

# Bubble Sort – how to work

89    45    68    90    29    34    17

89 ↔ 45    68    90    29    34    17

45    89 ↔ 68    90    29    34    17

45    68    89 ↔ 90 ↔ 29    34    17

45    68    89    29    90 ↔ 34    17

45    68    89    29    34    90 ↔ 17

45    68    89    29    34    17 | **90**

45 ↔ 68 ↔ 89 ↔ 29    34    17 | **90**

45    68    29    89 ↔ 34    17 | **90**

45    68    29    34    89 ↔ 17 | **90**

45    68    29    34    17 | **89    90**

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# String Matching

- Given a string of  $n$  characters called the **text** and a string of  $m$  characters ( $m \leq n$ ) called the **pattern**, find a substring of the text that matches the pattern.
- A brute force algorithm for string matching:
  - Align the pattern against the first  $m$  characters of the text.
  - Start matching the corresponding pair of characters from left to right until either all the  $m$  pairs are match.
  - Or if the missing pair is found, the pattern is shifted one position to the right and character comparisons are resumed, starting again from the 1<sup>st</sup> character.

# String Matching – how to work

NOBODY\_NOTICED\_HIM

Text

**NOT**

Pattern

N

N

.....

NOT

# Sequential Search

- Given an array  $A.1, \dots, A.N$  and target value  $V$ .
- Search routine should return an index of  $V$  in  $A$ , if  $V$  is present in the array, and  $N+1$  otherwise.

Given array – 12 24 35 42 46 49 50 82 99

Target value – 52

Return value : 10

Given array – 12 24 35 42 46 49 50 82 99

Target value – 42

Return value : 4



# Exhaustive Search

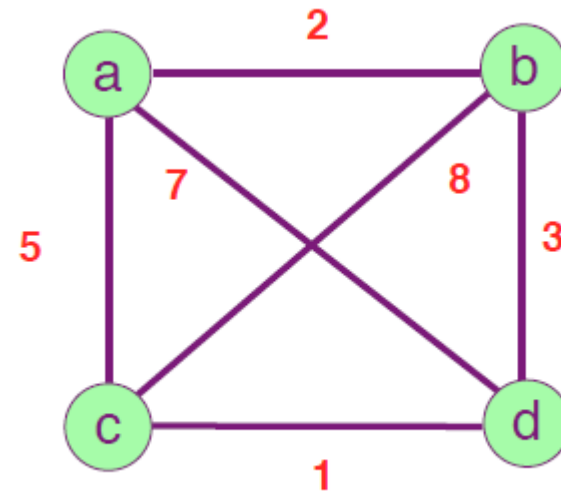
- Exhaustive search is a brute-force approach to combinatorial problems.
  - It suggests generating each and every combinatorial object of the problem, selecting those of them of that satisfy the problem's constraints and then finding a desired object.
  - Impractical for all but applicable to very small instances of problems.
- Examples:
  - Traveling salesman problem
    - Finding the shortest tour through a given set of  $n$  cities that visits each city exactly once before returning to the city where it started.
  - Knapsack problem
    - Finding the most valuable list of out-of  $n$  items that fit into the knapsack.
  - Job Assignment problem
    - Finding an assignment of  $n$  people to execute  $n$  jobs with the smallest total cost.

# Traveling Salesman Problem

- Given : A complete, weighted graph
- Find : A Hamilton circuit of minimum weight
- ✘ A circuit that starts at a vertex of a graph, passes through every vertex exactly once, and returns to the starting vertex is called a HAMILTON CIRCUIT.
- Algorithm
  - List all Hamilton circuits
  - Compute the sum of weights (total weight) for each circuit
  - Choose the one with the smallest total weight

# TSP – how to work

Tour	Length
a b c d a	$2+8+1+7 = 18$
a b d c a	$2+3+1+5 = \mathbf{11}$
a c b d a	23
a c d b a	<b>11</b>
a d b c a	23
a d c b a	18



More than one optimal solutions.

# Hamilton Circuits

- Step 1:*** Choose a starting point.
- Step 2:*** List all the Hamilton circuits with that starting point.
- Step 3:*** Find the total weight of each circuit.
- Step 4:*** Choose a Hamilton circuit with the smallest total weight.

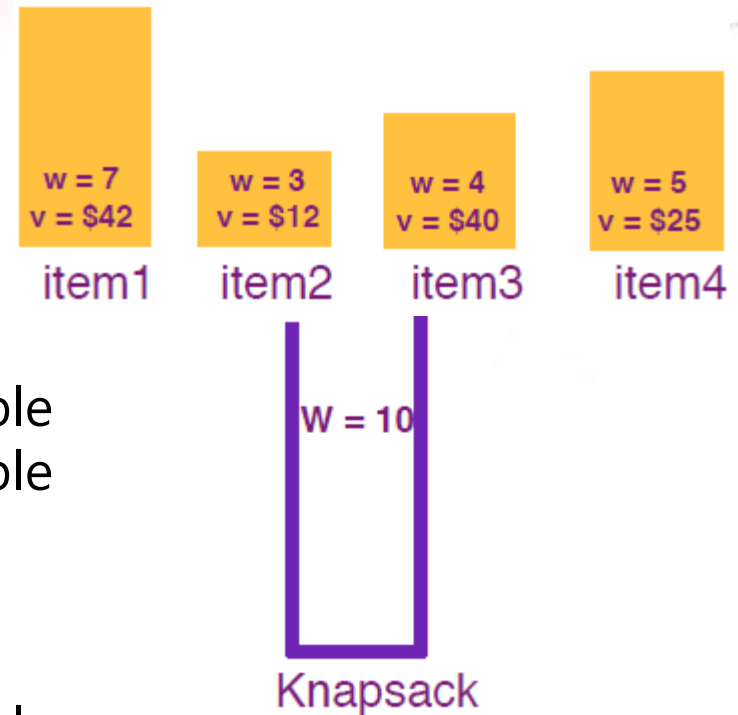
# Knapsack Problem

- Given  $n$  items of known weights  $w_1, \dots, w_n$  and values  $v_1, \dots, v_n$  and a knapsack of capacity  $W$ , find the most valuable subset of the items that fit into the knapsack.
- Consider all the subsets of the set of  $n$  items given,
  - Computing the total weight of each subset in order to identify feasible subsets (the ones with the total not exceeding the knapsack's capacity).
  - Finding a subset of the largest value among them.



# Knapsack – how to work

Subset	Total weight	Total value
Null	0	\$0
{1}	7	\$42
{2}	3	\$12
{3}	4	\$40
{4}	5	\$25
{1,2}	10	\$36
{1,3}	11	not feasible
{1,4}	12	not feasible
{2,3}	7	\$52
{2,4}	8	\$37
{3,4}	9	\$65
{1,2,3}	14	not feasible
{1,2,4}	15	not feasible
{1,3,4}	16	not feasible
{2,3,4}	12	not feasible
{1,2,3,4}	19	not feasible



# Job Assignment Problem

- There are  $n$  people who need to be assigned to  $n$  jobs, one person per job. The cost of assigning person  $i$  to job  $j$  is  $C[i,j]$ . Find an assignment that minimizes the total cost.
- Select one element in each row so that all selected elements are in different columns and the total sum of the selected elements is the smallest possible.

# Job Assignment Problem

	Job 1	Job 2	Job 3	Job 4
Person 1	<b>9</b>	<b>2</b>	<b>7</b>	<b>8</b>
Person 2	<b>6</b>	<b>4</b>	<b>3</b>	<b>7</b>
Person 3	<b>5</b>	<b>8</b>	<b>1</b>	<b>8</b>
Person 4	<b>7</b>	<b>6</b>	<b>9</b>	<b>4</b>

$$C = \begin{bmatrix} 9 & 2 & 7 & 8 \\ 6 & 4 & 3 & 7 \\ 5 & 8 & 1 & 8 \\ 7 & 6 & 9 & 4 \end{bmatrix}$$

Assignment (col.#s)

1, 2, 3, 4

1, 2, 4, 3

1, 3, 2, 4

1, 3, 4, 2

1, 4, 2, 3

1, 4, 3, 2

Total Cost

$$9+4+1+4=18$$

$$9+4+8+9=30$$

$$9+3+8+4=24$$

$$9+3+8+6=26$$

$$9+7+8+9=33$$

$$9+7+1+6=23$$

etc.

# Strengths and Weaknesses

- Strengths
  - wide applicability
  - simplicity
  - yields reasonable algorithms for some important problems (e.g., matrix multiplication, sorting, searching, string matching)
- Weaknesses
  - rarely yields efficient algorithms
  - some brute-force algorithms are unacceptably slow
  - not as constructive/creative as some other design techniques

# Q & A